## Docmosis v2.0 Release Notes

## New Features

Ch	ange	
Ta	ble row repetition i	s now defined differently
	11((1 1 2)	
Ine	e old "by-example"	mechanism of copying rows is <b>no longer supported</b> .
Old	ł way:	
	ID	Value
	«id»	«value»
	«id»	«value»
Nev	w way:	Value
	ID	Value
	uid>	////alua
	«iu»	«value»
	5. Alternating row calling to a reno template). The colour and the b provides the od	lerer (though renderers can still override the colour from the background colour of the $rr_{\rm row}$ is used as the even row background colour of each template row within the $rr_{\rm section}$ d row colour. Thus the above example would render like this
	ID	Value
	1	Alpha
	2	Beta
	3	Gamma
	4	Omega
	More complex c	olour alternation is possible:
	ID	Value
	«rr_idRow»	
	«id»	«value»
	«id»	«value»
1	«er idRow»	

	ID	Value
	1	Alpha
	1	Alpha
	2	Beta
	2	Beta
	3	Gamma
	3	Gamma
	4	Omega
	4	Omega
	bounding the rov	vs as desired: <i>Value</i>
	«rr idRow»	
	«id»	«value»
	«er idRow»	
	Would render as:	Valua
	1	Δ lpha
	2	Beta
	3	Gamma
	4	Omega
	5. Colouring and be	order styling are prototyped from the <i>rr</i> and <i>er</i> rows a
Tei	mentioned. This same number of done. mplate Errors to Ou	applies on a cell-by-cell basis if your content rows have columns as the <i>rr</i> _ and <i>er</i> _ rows. Otherwise a best-effor <b>tput Document</b>

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	table is all abo	out Greek		
1115				
	1D bb1	Value	Alpha	
	bb2	<	<value>&gt;<sup>Docmosis 1</sup></value>	
	bb«id»		«value»	
Docmo faik Cat end Su	sis1 Renderer ed with the exception use:java lang Arrayl er(line:320) ggestion 1:Look into	" (name="alt1", class=clas n below when trying to re Index OutOfBounds Excep the Renderer implement	rs com.docmosis.example.Ex nder field "value". tion:com.docmosis.example ation for the cause	xamp leRe nders \$3) .Example Renders \$3 x
This behavio	ur can be disa opulator.erro	bled by setting th <i>r.fatal=true</i>	e property:	
NOTE: this r	new style of en	rror handling is p	ogressive and curr	ently doesn't cover all
Conditional	Table Colum	nns		
	and variables	are now cunnorta	d so this form of ac	nditioning expression
are also supp in notes to fo	and variables orted "cc_{\$c ollow]. For ex	are now supporte columns<2}". [V ample:	d so this form of co ariables and expres	onditioning expression ssions are introduced
are also supp in notes to fo	and variables orted "cc_{\$c ollow]. For ex	are now supporte columns<2}". [V ample: <i>Value «cc_hasP</i>	d so this form of co ariables and expres <b>rimeValues</b> »	onditioning expression sions are introduced
are also supp in notes to fo	and variables orted "cc_{\$c ollow]. For ex	are now supporte columns<2}". [V ample: <i>Value «cc_hasP</i>	d so this form of co ariables and expres <b>rimeValues</b> »	onditioning expression ssions are introduced
are also supp in notes to fo	and variables orted "cc_{\$c ollow]. For ex idRow»	are now supporte columns<2}". [V ample: <i>Value «cc_hasP</i>	d so this form of co ariables and expres <b>rimeValues</b> » «value»	onditioning expression sions are introduced
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are also supp in notes to for <i>ID</i> «rr «er The second or "hasPrimeVa columns ther the left as the This goes a b columns. Co	and variables orted "cc_{\$c ollow]. For ex idRow» idRow» column will be alues". Note t a this will alway e marked column oit further in the taken of taken of the taken of t	are now supporte columns<2}". [V ample: <i>Value «cc_hasP</i> e removed if the d hat if your table h ays remove the co mn. hat the condition v ole:	d so this form of co ariables and expres rimeValues» «value» ata source returns a as rows with differ olumn at the same r	a false for relative location from
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are also supp in notes to for <i>ID</i> «rr «er The second or "hasPrimeVa columns then the left as the Columns. Co	and variables orted "cc_{\$c ollow]. For ex idRow» «id» idRow» column will be alues". Note t a this will alwa e marked colum of further in the onsider this tab	are now supporte columns<2}". [V ample: Value «cc_hasP e removed if the d hat if your table h ays remove the co mn. hat the condition v ole: Value «cc_hasP Value «cc_hasP	d so this form of co ariables and expres rimeValues» «value» ata source returns a as rows with differ olumn at the same r will also be applied rimeValues» 2» «value	a false for relative location from to sub (or covered)
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	such, these columns will also be removed if the value column is removed.		
4	Variables		
	Simple variables can be defined in the template and be referenced later. Defining a variable requires a merge field of the form <i>\$var=term</i> . For example:		
	«\$a=hotels.floors[1]»		
	defines a variable named "a" which holds reference to the value for the second floor of a hotel. This variable may then be referenced later, for example this will get the name of the floor:		
	«\$a.floorName»		
	and this will repeat of	over all the rooms in the floor:	
	«rs_\$a.rooms» «roomNumber» «es_\$a.rooms»		
	NOTE: support for l	literal values and expressions when setting variables is pend	ding.
5	Built in variables There are some built-in variables to help with referencing data within or without the "current" context of data.		
	Variable	Description	
	Scurrent	Current data context Allowed unnamed types of	
	Sthis	data to be referenced such as arrays or collections of	
	ψtins	Iava primitives For example Scurrent[0] might	
		reference the first string in an array of strings	
	\$top	The root of the data – regardless of where in the	
	\$root	population phase the render process is, the data can be absolutely referenced.	
	\$parent	The data one level up from the current population context. For example, in a table row that is populating rooms of a floor in a repeating fashion, the \$parent variable would refer to the parent of the rooms container, which is likely to be the current	
	\$idx	10or.         Current population index, that is, index into the	
		current data set.	
	\$itemnum	Current item number which is the position within the current data set where the starting position is 1.	

	\$size	The size of the current data set.		
6	Expression Support			
	Expressions can now	be used for conditional sections and conditional colum	ns. The	
	format of an expressi	on is:		
	«cs_{expr}» or «cc_{expr}»			
	where supported expressions are:			
	Expression	Description		
	{a <b}< th=""><th>Lookup data "b" less than lookup of data "a"</th><th></th></b}<>	Lookup data "b" less than lookup of data "a"		
	{\$a <b}< th=""><th>Lookup variable "a" less than lookup of data "b"</th><th></th></b}<>	Lookup variable "a" less than lookup of data "b"		
	{a='123'}	Lookup data "a" equals String literal "123"		
	{b!=123}	Lookup data "b" is not equal to Number literal 123		
	{!c}	Lookup data "c" and logically negate it		
	{!\$a}	Lookup variable "a" and logically negate it		
	Examples are:			
	«cc_{\$columns<10}»			
	«cs_{roomNumber<]	01}»		
7	Enhanced Range Sn	ecification		
'	Ennanced Range Sp			
	Ranges can now be si	pecified so that parts of arrays or collections of data ma	v be	
	referenced. For exam	iple:	5	
		L		
	Field	Description		
	«rooms[0]»	First room	7	
	«rooms[F]»	First room	7	
	«rooms[L]»	Last room	7	
	«rooms[L3]	» First 3 rooms	7	
	«rooms[F2,]	L2]» First 2 rooms and last 2 rooms		
	«rooms[1-3	>>> 2nd to 4 <sup>th</sup> room		
	«rooms[1,3,	L2]» 2 <sup>nd</sup> , 4 <sup>th</sup> and last 2 rooms		
	«rooms[1-L	2]» 2 <sup>nd</sup> to second last rooms		
	«rooms[*]»	All rooms		
8	Automatic Template	e Registration		
	Dethe con 1		1	
	raths can now be con	figured to be automatically monitored for templates an	a	
	l template changes. If	ie properties:		
	docmosis tomplate m	particer sourcement = path 1 meth?		
	docmosis template m	nuor.sourcepain-pain1,pain2		
	$\perp aocmosis.template.monitor.perioa=3$			

	if set will automatically watch for templates and register them according to their file name and relative location.
	NOTE: templates in jar files are also supported.
9	Simplified Data Provision
	All collecting of data to provide to Docmosis should now be performed via the DataProviderBuilder class. This class allows all forms of data to be attached, without needing to know anything about the underlying implementations.
	So, to provide data from some a combination of Java objects, a few strings of data and an SQL query, the code might look like this:
	DataProviderBuilder dpb = new DataProviderBuilder(); dpb.addJavaObject(myFriend);
	dpb.addJavaObject(myUncle); dpb.add("reportTitle", "Associate Report");
	dpb.add("footerLabel", "Associate Report as at " + myDateStr); dpb.addSQL(myResultSet);
	 DataProvider dataProvider = dpb.getDataProvider();
10	Simplified Template Identification
	All types of TemplateIdentifier are now represented by the TemplateIdentifier class. All other classes have been remove. The same applies to TemplateContext which now assumes all contexts can be defined by a "/" separated "path".
	Templates can now be identified with or without a Context and if no Context is provided, the template is assumed to be at the root of the template store.
	So, to identify a template, we now do:
	TemplateIdentifier tid = new TemplateIdentifier("myTemplate"); Or
	TemplateIdentifier tid = new TemplateIdentifier("myTemplate", "project1/release1");
11	Simplified Rendering

	File template = new File("myTemplate.doc"); File result = new File("output/myDocument.pdf"); DataProvider data = new DataProviderBuilder().getDataProvider(); DocumentProcessor.renderDoc(template, result, data);	
12	Improved Core Performance	
	<ul> <li>The core engine has had several improvements to reduce processing effort during document production including:</li> <li>A new cache that holds ready-to-render templates in memory (default configuration is 5Mb)</li> <li>Complex data lookup structures are pre-computed in the template registration phase.</li> </ul>	

## **API Changes**

The following API changes should be noted

Class / Interface	Change
ConversionInstruction	Interface changed into Class
SimpleConversionInstruction	Replaced with ConversionInstruction
FileTemplateContext	Replaced with TemplateContext
FileTemplateIdentifier	Replaced with TemplateIdentifier
TemplateContext	Interface changed to class
TemplateIdentifier	Interface changed to class
FieldDetails	Package change
RenderedField	Package change. Changed from Interface to class
SimpleRenderedField	Replaced by RenderedField
FieldRenderer	Package change

## Bug Fixes / Technical Enhancements

#	Change
1	Fixed a small memory leak that can occur when generating over 50000 documents
	without restarting a Docmosis converter.
2	Fixed an unhelpful error that occurs when renderers are not used correctly.
3	Fixed an unhelpful error that occurs when reflective data providers are not used
	correctly.
4	SQL Data Providers now handle all the different SQL data types.
5	Booleans and Strings can be treated interchangeably. Affects only String and File
	based data provision (not Reflective or SQL data).
6	Increased reliability of start-end tag matching despite styling in the document.